RBGL

April 19, 2010

astarSearch

Compute astarSearch for a graph

Description

Compute astarSearch for a graph

Usage

```
astarSearch(g)
```

Arguments

g

an instance of the graph class

Details

NOT IMPLEMENTED YET. TO BE FILLED IN

Author(s)

Li Long < li.long@isb-sib.ch>

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

The Boost Graph Library: User Guide and Reference Manual; by Jeremy G. Siek, Lie-Quan Lee, and Andrew Lumsdaine; (Addison-Wesley, Pearson Education Inc., 2002), xxiv+321pp. ISBN 0-201-72914-8

```
con <- file(system.file("XML/dijkex.gxl",package="RBGL"), open="r")
coex <- fromGXL(con)
close(con)
astarSearch(coex)</pre>
```

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bandwidth

Compute bandwidth for an undirected graph

Description

Compute bandwidth for an undirected graph

Usage

```
bandwidth(g)
```

Arguments

g

an instance of the graph class with edgemode "undirected"

Details

The bandwidth of an undirected graph G=(V, E) is the maximum distance between two adjacent vertices. See documentation on bandwidth in Boost Graph Library for more details.

Value

```
bandwidth the bandwidth of the given graph
```

Author(s)

Li Long < li.long@isb-sib.ch>

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

The Boost Graph Library: User Guide and Reference Manual; by Jeremy G. Siek, Lie-Quan Lee, and Andrew Lumsdaine; (Addison-Wesley, Pearson Education Inc., 2002), xxiv+321pp. ISBN 0-201-72914-8

```
con <- file(system.file("XML/dijkex.gxl",package="RBGL"), open="r")
coex <- fromGXL(con)
close(con)
coex <- ugraph(coex)
bandwidth(coex)</pre>
```

```
betweenness.centrality.clustering
```

Graph clustering based on edge betweenness centrality

Description

Graph clustering based on edge betweenness centrality

Usage

```
betweenness.centrality.clustering(g, threshold = -1, normalize = T)
```

Arguments

g an instance of the graph class with edgemode "undirected"

threshold to terminate clustering process

normalize boolean, when TRUE, the edge betweenness centrality is scaled by 2/((n-

1) (n-2)) where n is the number of vertices in g; when FALSE, the edge

betweenness centrality is the absolute value

Details

To implement graph clustering based on edge betweenness centrality.

The algorithm is iterative, at each step it computes the edge betweenness centrality and removes the edge with maximum betweenness centrality when it is above the given threshold. When the maximum betweenness centrality falls below the threshold, the algorithm terminates.

See documentation on Clustering algorithms in Boost Graph Library for details.

Value

A list of

```
no.of.edges number of remaining edges after removal
```

edges remaining edges edge.betweenness.centrality

betweenness centrality of remaining edges

Author(s)

Li Long <li.long@isb-sib.ch>

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

The Boost Graph Library: User Guide and Reference Manual; by Jeremy G. Siek, Lie-Quan Lee, and Andrew Lumsdaine; (Addison-Wesley, Pearson Education Inc., 2002), xxiv+321pp. ISBN 0-201-72914-8

See Also

```
brandes.betweenness.centrality
```

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Examples

```
con <- file(system.file("XML/conn.gxl",package="RBGL"))
coex <- fromGXL(con)
close(con)
coex <- ugraph(coex)
betweenness.centrality.clustering(coex, 0.5, TRUE)</pre>
```

bellman.ford.sp

Bellman-Ford shortest paths using boost C++

Description

Algorithm for the single-source shortest paths problem for a graph with both positive and negative edge weights.

Usage

```
bellman.ford.sp(g, start=nodes(g)[1])
```

Arguments

g instance of class graph

start character: node name for start of path

Details

This function interfaces to the Boost graph library C++ routines for Bellman-Ford shortest paths. Choose the appropriate algorithm to calculate the shortest path carefully based on the properties of the given graph. See documentation on Bellman-Ford algorithm in Boost Graph Library for more details.

Value

A list with elements:

all edges minimized

true if all edges are minimized, false otherwise.

distance The vector of distances from start to each node of g; includes Inf when

there is no path from start.

penult A vector of indices (in nodes (g)) of predecessors corresponding to each node

on the path from that node back to start

. For example, if the element one of this vector has value 10, that means that the predecessor of node 1 is node 10. The next predecessor is found by examining penult [10].

start The start node that was supplied in the call to bellman.ford.sp.

Author(s)

Li Long <li.long@isb-sib.ch>

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

The Boost Graph Library: User Guide and Reference Manual; by Jeremy G. Siek, Lie-Quan Lee, and Andrew Lumsdaine; (Addison-Wesley, Pearson Education Inc., 2002), xxiv+321pp. ISBN 0-201-72914-8

See Also

```
dag.sp, dijkstra.sp, johnson.all.pairs.sp, sp.between
```

Examples

```
con <- file(system.file("XML/conn2.gxl",package="RBGL"), open="r")
dd <- fromGXL(con)
close(con)
bellman.ford.sp(dd)
bellman.ford.sp(dd,nodes(dd)[2])</pre>
```

brandes.betweenness.centrality

Compute betweenness centrality for an undirected graph

Description

Compute betweenness centrality for an undirected graph

Usage

```
brandes.betweenness.centrality(g)
```

Arguments

g

an instance of the graph class with edgemode "undirected"

Details

Brandes.betweenness.centrality computes the betweenness centrality of each vertex or each edge in the graph, using an algorithm by U. Brandes.

Betweenness centrality of a vertex v is calculated as follows: $N_st(v) = no$. of shortest paths from s to t that pass through v, $N_st = no$. of shortest paths from s to t, betweenness centrality of $v = sum(N_st(v)/N_st)$ for all vertices s!=v!=t.

Betweenness centrality of an edge is calculated similarly.

The relative betweenness centrality for a vertex is to scale the betweenness centrality of the given vertex by 2/(n**2 - 3n + 2) where n is the no. of vertices in the graph.

Central point dominance measures the maximum betweenness of any vertex in the graph.

See documentation on brandes betweenness centrality in Boost Graph Library for more details.

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Value

```
A list of
```

```
betweenness.centrality.vertices
betweenness.centrality.edges
betweenness.centrality of each edge
relative.betweenness.centrality.vertices
relative betweenness centrality of each vertex
dominance maximum betweenness of any point in the graph
```

Author(s)

Li Long <li.long@isb-sib.ch>

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

The Boost Graph Library: User Guide and Reference Manual; by Jeremy G. Siek, Lie-Quan Lee, and Andrew Lumsdaine; (Addison-Wesley, Pearson Education Inc., 2002), xxiv+321pp. ISBN 0-201-72914-8

See Also

```
betweenness.centrality.clustering
```

Examples

```
con <- file(system.file("XML/conn.gxl",package="RBGL"), open="r")
coex <- fromGXL(con)
close(con)
coex <- ugraph(coex)
brandes.betweenness.centrality(coex)</pre>
```

bfs

Breadth and Depth-first search

Description

These functions return information on graph traversal by breadth and depth first search using routines from the BOOST library.

Usage

```
bfs(object, node, checkConn=TRUE)
dfs(object, node, checkConn=TRUE)
```

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Arguments

object instance of class graph from Bioconductor graph class

node node name where search starts; defaults to the node in first position in the node

vector.

checkConn logical for backwards compatibility; this parameter has no effect as of RBGL

1.7.9 and will be removed in future versions.

Details

These two functions are interfaces to the BOOST graph library functions for breadth first and depth first search. Both methods handle unconnected graphs by applying the algorithms over the connected components.

Cormen et al note (p 542) that 'results of depth-first search may depend upon the order in which the vertices are examined ... These different visitation orders tend not to cause problems in practice, as any DFS result can usually be used effectively, with essentially equivalent results'.

Value

For bfs a vector of node indices in order of BFS visit.

For dfs a list of two vectors of nodes, with elements discover (order of DFS discovery), and finish (order of DFS completion).

Author(s)

VJ Carey <stvjc@channing.harvard.edu>

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

The Boost Graph Library: User Guide and Reference Manual; by Jeremy G. Siek, Lie-Quan Lee, and Andrew Lumsdaine; (Addison-Wesley, Pearson Education Inc., 2002), xxiv+321pp. ISBN 0-201-72914-8

```
con1 <- file(system.file("XML/bfsex.gxl",package="RBGL"), open="r")
dd <- fromGXL(con1)
close(con1)

bfs(dd, "r")
bfs(dd, "s")

con2 <- file(system.file("XML/dfsex.gxl",package="RBGL"), open="r")
dd2 <- fromGXL(con2)
close(con2)

dfs(dd2, "u")</pre>
```

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biConnComp

Compute biconnected components for a graph

Description

Compute biconnected components for a graph

Usage

```
biConnComp(g)
articulationPoints(g)
```

Arguments

g

an instance of the graph class

Details

A biconnected graph is a connected graph that remains connected when any one of its vertices, and all the edges incident on this vertex, is removed and the graph remains connected. A biconnected component of a graph is a subgraph which is biconnected. An integer label is assigned to each edge to indicate which biconnected component it's in.

A vertex in a graph is called an articulation point if removing it increases the number of connected components.

See the documentation for the Boost Graph Library for more details.

Value

For biConnComp: a vector whose length is no. of biconnected components, each entry is a list of nodes that are on the same biconnected components.

For articulationPoints: a vector of articulation points in the graph.

Author(s)

Li Long <li.long@isb-sib.ch>

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

The Boost Graph Library: User Guide and Reference Manual; by Jeremy G. Siek, Lie-Quan Lee, and Andrew Lumsdaine; (Addison-Wesley, Pearson Education Inc., 2002), xxiv+321pp. ISBN 0-201-72914-8

```
con <- file(system.file("XML/conn.gxl",package="RBGL"), open="r")
coex <- fromGXL(con)
close(con)
biConnComp(coex)
articulationPoints(coex)</pre>
```

 ${\tt boyerMyrvoldPlanarityTest}$

boyer Myrvold Planarity Test

Description

boyerMyrvoldPlanarityTest description

Usage

boyerMyrvoldPlanarityTest(g)

Arguments

g instance of class graphNEL from Bioconductor graph class

Author(s)

Li Long <li.long@isb-sib.ch>

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

The Boost Graph Library: User Guide and Reference Manual; by Jeremy G. Siek, Lie-Quan Lee, and Andrew Lumsdaine; (Addison-Wesley, Pearson Education Inc., 2002), xxiv+321pp. ISBN 0-201-72914-8

Description

chrobakPayneStraightLineDrawing description

Usage

chrobakPayneStraightLineDrawing(g)

Arguments

g instance of class graphNEL from Bioconductor graph class

Author(s)

Li Long <li.long@isb-sib.ch>

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References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

The Boost Graph Library: User Guide and Reference Manual; by Jeremy G. Siek, Lie-Quan Lee, and Andrew Lumsdaine; (Addison-Wesley, Pearson Education Inc., 2002), xxiv+321pp. ISBN 0-201-72914-8

Examples

```
V <- LETTERS[1:7]
g <- new("graphNEL", nodes=V, edgemode="undirected")</pre>
q \leftarrow addEdge(V[1+0], V[1+1], q)
g \leftarrow addEdge(V[1+1], V[2+1], g)
g \leftarrow addEdge(V[1+2], V[3+1], g)
g \leftarrow addEdge(V[1+3], V[0+1], g)
g \leftarrow addEdge(V[1+3], V[4+1], g)
g \leftarrow addEdge(V[1+4], V[5+1], g)
g \leftarrow addEdge(V[1+5], V[6+1], g)
g \leftarrow addEdge(V[1+6], V[3+1], g)
g \leftarrow addEdge(V[1+0], V[4+1], g)
g \leftarrow addEdge(V[1+1], V[3+1], g)
g \leftarrow addEdge(V[1+3], V[5+1], g)
g \leftarrow addEdge(V[1+2], V[6+1], g)
g \leftarrow addEdge(V[1+1], V[4+1], g)
g \leftarrow addEdge(V[1+1], V[5+1], g)
g \leftarrow addEdge(V[1+1], V[6+1], g)
x3 <- chrobakPayneStraightLineDrawing(g)
x3
```

clusteringCoefAppr Approximate clustering coefficient for an undirected graph

Description

Approximate clustering coefficient for an undirected graph

Usage

```
clusteringCoefAppr(g, k=length(nodes(g)), Weighted=FALSE, vW=degree(g))
```

Arguments

```
g an instance of the graph class

Weighted calculate weighted clustering coefficient or not

vW vertex weights to use when calculating weighted clustering coefficient

k parameter controls total expected runtime
```

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Details

It is quite expensive to compute cluster coefficient and transitivity exactly for a large graph by computing the number of triangles in the graph. Instead, clusteringCoefAppr samples triples with appropriate probability, returns the ratio between the number of existing edges and the number of samples.

MORE ABOUT CHOICE OF K.

See reference for more details.

Value

Approximated clustering coefficient for graph q.

Author(s)

Li Long <li.long@isb-sib.ch>

References

Approximating Clustering Coefficient and Transitivity, T. Schank, D. Wagner, Journal of Graph Algorithms and Applications, Vol. 9, No. 2 (2005).

See Also

clusteringCoef, transitivity, graphGenerator

Examples

```
con <- file(system.file("XML/conn.gxl",package="RBGL"))
g <- fromGXL(con)
close(con)

k = length(nodes(g))
cc <- clusteringCoefAppr(g, k)
ccwl <- clusteringCoefAppr(g, k, Weighted=TRUE)
vW <- c(1, 1, 1, 1, 1, 1, 1)
ccw2 <- clusteringCoefAppr(g, k, Weighted=TRUE, vW)</pre>
```

clusteringCoef

Calculate clustering coefficient for an undirected graph

Description

Calculate clustering coefficient for an undirected graph

Usage

```
clusteringCoef(g, Weighted=FALSE, vW=degree(g))
```

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Arguments

g	an instance of the graph class
Weighted	calculate weighted clustering coefficient or not
√M	vertex weights to use when calculating weighted clustering coefficient

Details

For an undirected graph G, let delta(v) be the number of triangles with v as a node, let tau(v) be the number of triples, i.e., paths of length 2 with v as the center node.

Let V' be the set of nodes with degree at least 2.

Define clustering coefficient for v, c(v) = (delta(v) / tau(v)).

Define clustering coefficient for G, C(G) = sum(c(v)) / |V'|, for all \forall in V'.

Define weighted clustering coefficient for g, Cw(G) = sum(w(v) * c(v)) / sum(w(v)), for all v in V

Value

Clustering coefficient for graph G.

Author(s)

```
LiLong <li.long@isb-sib.ch>
```

References

Approximating Clustering Coefficient and Transitivity, T. Schank, D. Wagner, Journal of Graph Algorithms and Applications, Vol. 9, No. 2 (2005).

See Also

clusteringCoefAppr, transitivity, graphGenerator

```
con <- file(system.file("XML/conn.gxl",package="RBGL"))
g <- fromGXL(con)
close(con)
cc <- clusteringCoef(g)
ccw1 <- clusteringCoef(g, Weighted=TRUE)
vW <- c(1, 1, 1, 1, 1, 1, 1)
ccw2 <- clusteringCoef(g, Weighted=TRUE, vW)</pre>
```

```
sequential.vertex.coloring
```

Compute a vertex coloring for a graph

Description

Compute vertex coloring for a graph

Usage

```
sequential.vertex.coloring(g)
```

Arguments

g

an instance of the graph class

Details

A vertex coloring for a graph is to assign a color for each vertex so that no two adjacent vertices are of the same color. We designate the colors as sequential integers: 1, 2,

For ordered vertices, v1, v2, ..., vn, for k = 1, 2, ..., n, this algorithm assigns vk to the smallest possible color. It does NOT guarantee to use minimum number of colors.

See documentations on these algorithms in Boost Graph Library for more details.

Value

```
no. of colors needed

how many colors to use to color the graph

colors of nodes

color label for each vertex
```

Author(s)

Li Long <li.long@isb-sib.ch>

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

The Boost Graph Library: User Guide and Reference Manual; by Jeremy G. Siek, Lie-Quan Lee, and Andrew Lumsdaine; (Addison-Wesley, Pearson Education Inc., 2002), xxiv+321pp. ISBN 0-201-72914-8

```
con <- file(system.file("XML/dijkex.gxl",package="RBGL"), open="r")
coex <- fromGXL(con)
close(con)
sequential.vertex.coloring(coex)</pre>
```

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connectedComp

Identify Connected Components in an Undirected Graph

Description

The connected components in an undirected graph are identified. If the graph is directed then the weakly connected components are identified.

Usage

```
connectedComp(g)
```

Arguments

g

graph with edgemode "undirected"

Details

Uses a depth first search approach to identifying all the connected components of an undirected graph. If the input, g, is a directed graph it is first transformed to an undirected graph (using ugraph).

Value

A list of length equal to the number of connected components in g. Each element of the list contains a vector of the node labels for the nodes that are connected.

Author(s)

Vince Carey <stvjc@channing.harvard.edu>

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

The Boost Graph Library: User Guide and Reference Manual; by Jeremy G. Siek, Lie-Quan Lee, and Andrew Lumsdaine; (Addison-Wesley, Pearson Education Inc., 2002), xxiv+321pp. ISBN 0-201-72914-8

See Also

connComp,strongComp, ugraph, same.component

```
con <- file(system.file("GXL/kmstEx.gxl",package="graph"), open="r")
km <- fromGXL(con)
close(con)
km <- graph::addNode(c("F","G","H"), km)
km <- addEdge("G", "H", km, 1)
km <- addEdge("H", "G", km, 1)
ukm <- ugraph(km)
ukm
edges(ukm)
connectedComp(ukm)</pre>
```

dag.sp 15

dag.sp DAG shortest paths using boost (

Description

Algorithm for the single-source shortest-paths problem on a weighted, directed acyclic graph (DAG)

Usage

```
dag.sp(g, start=nodes(g)[1])
```

Arguments

g instance of class graph start source node for start of paths

Details

These functions are interfaces to the Boost graph library C++ routines for single-source shortest-paths on a weighted directed acyclic graph. Choose appropriate shortest-path algorithms carefully based on the properties of the input graph. See documentation in Boost Graph Library for more details.

Value

A list with elements:

distance The vector of distances from start to each node of g; includes Inf when

there is no path from start.

penult A vector of indices (in nodes (g)) of predecessors corresponding to each node

on the path from that node back to start. For example, if the element one of this vector has value 10, that means that the predecessor of node 1 is node 10.

The next predecessor is found by examining penult [10].

start The start node that was supplied in the call to dag.sp.

Author(s)

Li Long li.long@isb-sib.ch>

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

The Boost Graph Library: User Guide and Reference Manual; by Jeremy G. Siek, Lie-Quan Lee, and Andrew Lumsdaine; (Addison-Wesley, Pearson Education Inc., 2002), xxiv+321pp. ISBN 0-201-72914-8

See Also

```
bellman.ford.sp, dijkstra.sp, johnson.all.pairs.sp, sp.between
```

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Examples

```
con <- file(system.file("XML/conn2.gxl",package="RBGL"), open="r")
dd <- fromGXL(con)
close(con)
dag.sp(dd)
dag.sp(dd,nodes(dd)[2])</pre>
```

dijkstra.sp

Dijkstra's shortest paths using boost C++

Description

dijkstra's shortest paths

Usage

```
dijkstra.sp(g, start=nodes(g)[1], eW=unlist(edgeWeights(g)))
```

Arguments

g instance of class graph

start character: node name for start of path

eW numeric: edge weights.

Details

These functions are interfaces to the Boost graph library C++ routines for Dijkstra's shortest paths.

For some graph subclasses, computing the edge weights can be expensive. If you are calling dijkstra.sp in a loop, you can pass the edge weights explicitly to avoid the edge weight creation cost.

Value

A list with elements:

distance The vector of distances from start to each node of g; includes Inf when

there is no path from start.

penult A vector of indices (in nodes (g)) of predecessors corresponding to each node

on the path from that node back to start

. For example, if the element one of this vector has value 10, that means that the predecessor of node 1 is node 10. The next predecessor is found by examining penult [10].

start The start node that was supplied in the call to dijkstra.sp.

Author(s)

VJ Carey <stvjc@channing.harvard.edu>

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References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

The Boost Graph Library: User Guide and Reference Manual; by Jeremy G. Siek, Lie-Quan Lee, and Andrew Lumsdaine; (Addison-Wesley, Pearson Education Inc., 2002), xxiv+321pp. ISBN 0-201-72914-8

See Also

```
bellman.ford.sp, dag.sp, johnson.all.pairs.sp, sp.between
```

Examples

```
con1 <- file(system.file("XML/dijkex.gxl",package="RBGL"), open="r")
dd <- fromGXL(con1)
close(con1)
dijkstra.sp(dd)
dijkstra.sp(dd,nodes(dd)[2])

con2 <- file(system.file("XML/ospf.gxl",package="RBGL"), open="r")
ospf <- fromGXL(con2)
close(con2)
dijkstra.sp(ospf,nodes(ospf)[6])</pre>
```

dominatorTree

Compute dominator tree from a vertex in a directed graph

Description

Compute dominator tree from a vertex in a directed graph

Usage

```
dominatorTree(g, start=nodes(g)[1])
lengauerTarjanDominatorTree(g, start=nodes(g)[1])
```

Arguments

```
g a directed graph, one instance of the graph class start a vertex in graph g
```

Details

As stated in documentation on Lengauer Tarjan dominator tree in Boost Graph Library:

A vertex u dominates a vertex v, if every path of directed graph from the entry to v must go through u.

This function builds the dominator tree for a directed graph.

Value

Output is a vector, giving each node its immediate dominator.

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Author(s)

Li Long <li.long@isb-sib.ch>

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

Examples

```
con1 <- file(system.file("XML/dominator.gxl",package="RBGL"), open="r")
g1 <- fromGXL(con1)
close(con1)
dominatorTree(g1)
lengauerTarjanDominatorTree(g1)</pre>
```

edgeConnectivity

computed edge connectivity and min disconnecting set for an undirected graph

Description

computed edge connectivity and min disconnecting set for an undirected graph

Usage

```
edgeConnectivity(g)
```

Arguments

g

an instance of the graph class with edgemode "undirected"

Details

Consider a graph G consisting of a single connected component. The edge connectivity of G is the minimum number of edges in G that can be cut to produce a graph with two (disconnected) components. The set of edges in this cut is called the minimum disconnecting set.

Value

A list:

connectivity the integer describing the number of edges that must be severed to obtain two components

minDisconSet a list (of length connectivity) of pairs of node names describing the edges

that need to be cut to obtain two components

Author(s)

Vince Carey <stvjc@channing.harvard.edu>

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

The Boost Graph Library: User Guide and Reference Manual; by Jeremy G. Siek, Lie-Quan Lee, and Andrew Lumsdaine; (Addison-Wesley, Pearson Education Inc., 2002), xxiv+321pp. ISBN 0-201-72914-8

See Also

```
minCut, edmonds.karp.max.flow, push.relabel.max.flow
```

Examples

```
con <- file(system.file("XML/conn.gxl",package="RBGL"), open="r")
coex <- fromGXL(con)
close(con)
edgeConnectivity(coex)</pre>
```

 ${\it edmonds Max Cardinality Matching} \\ {\it edmonds Max Cardinality Matching}$

Description

edmondsMaxCardinalityMatching description

Usage

```
edmondsMaxCardinalityMatching(g)
```

Arguments

~

instance of class graphNEL from Bioconductor graph class

Author(s)

Li Long < li.long@isb-sib.ch>

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

The Boost Graph Library: User Guide and Reference Manual; by Jeremy G. Siek, Lie-Quan Lee, and Andrew Lumsdaine; (Addison-Wesley, Pearson Education Inc., 2002), xxiv+321pp. ISBN 0-201-72914-8

Examples

```
V <- LETTERS[1:18]</pre>
g <- new("graphNEL", nodes=V, edgemode="undirected")</pre>
g \leftarrow addEdge(V[1+0], V[4+1], g);
g \leftarrow addEdge(V[1+1], V[5+1], g);
g \leftarrow addEdge(V[1+2], V[6+1], g);
g \leftarrow addEdge(V[1+3], V[7+1], g);
g <- addEdge(V[1+4], V[5+1], g);</pre>
g \leftarrow addEdge(V[1+6], V[7+1], g);
g \leftarrow addEdge(V[1+4], V[8+1], g);
g \leftarrow addEdge(V[1+5], V[9+1], g);
g <- addEdge(V[1+6], V[10+1], g);
g <- addEdge(V[1+7], V[11+1], g);</pre>
g <- addEdge(V[1+8], V[9+1], g);</pre>
g \leftarrow addEdge(V[1+10], V[11+1], g);
g \leftarrow addEdge(V[1+8], V[13+1], g);
g <- addEdge(V[1+9], V[14+1], g);</pre>
g <- addEdge(V[1+10], V[15+1], g);</pre>
g <- addEdge(V[1+11], V[16+1], g);</pre>
g \leftarrow addEdge(V[1+14], V[15+1], g);
x9 <- edmondsMaxCardinalityMatching(g)
x9
g \leftarrow addEdge(V[1+12], V[13+1], g);
g \leftarrow addEdge(V[1+16], V[17+1], g);
x10 <- edmondsMaxCardinalityMatching(g)</pre>
x10
```

edmondsOptimumBranching

edmondsOptimumBranching

Description

edmondsOptimumBranching description

Usage

```
edmondsOptimumBranching(g)
```

Arguments

J

instance of class graphNEL from Bioconductor graph class

Details

This is an implementation of Edmonds' algorithm to find optimum branching in a directed graph. See references for details.

extractPath 21

Author(s)

Li Long <li.long@isb-sib.ch>

References

See Edmonds' Algorithm on http://edmonds-alg.sourceforge.net.

Examples

```
V <- LETTERS[1:4]
g <- new("graphNEL", nodes=V, edgemode="directed")</pre>
g \leftarrow addEdge(V[1+0],V[1+1],g, 3)
g \leftarrow addEdge(V[1+0],V[2+1],g, 1.5)
g \leftarrow addEdge(V[1+0],V[3+1],g, 1.8)
g \leftarrow addEdge(V[1+1],V[2+1],g, 4.3)
g \leftarrow addEdge(V[1+2],V[3+1],g, 2.2)
x11 <- edmondsOptimumBranching(g)</pre>
x11
```

extractPath

convert a dijkstra.sp predecessor structure into the path joining two nodes

Description

determine a path between two nodes in a graph, using output of dijkstra.sp.

Usage

```
extractPath(s, f, pens)
```

Arguments

index of starting node in nodes vector of the graph from which pens was de-S rived f

index of ending node in nodes vector

predecessor index vector as returned in the preds component of dijkstra.sp pens output

Author(s)

Vince Carey <stvjc@channing.harvard.edu>

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

The Boost Graph Library: User Guide and Reference Manual; by Jeremy G. Siek, Lie-Quan Lee, and Andrew Lumsdaine; (Addison-Wesley, Pearson Education Inc., 2002), xxiv+321pp. ISBN 0-201-72914-8

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See Also

```
allShortestPaths
```

Examples

```
data(FileDep)
dd <- dijkstra.sp(FileDep)
extractPath(1,9,dd$pen)</pre>
```

FileDep

FileDep: a graphNEL object representing a file dependency dataset example in boost graph library

Description

FileDep: a graphNEL object representing a file dependency dataset example in boost graph library

Usage

```
#data(FileDep)
```

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

The Boost Graph Library: User Guide and Reference Manual; by Jeremy G. Siek, Lie-Quan Lee, and Andrew Lumsdaine; (Addison-Wesley, Pearson Education Inc., 2002), xxiv+321pp. ISBN 0-201-72914-8

```
# this is how the graph of data(FileDep) was obtained
library(graph)
fd <- file(system.file("XML/FileDep.gxl",package="RBGL"), open="r")
show(fromGXL(fd))
if (require(Rgraphviz))
{
   data(FileDep)
   plot(FileDep)
}
close(fd)</pre>
```

```
floyd.warshall.all.pairs.sp

compute shortest paths for all pairs of nodes
```

Description

compute shortest paths for all pairs of nodes

Usage

```
floyd.warshall.all.pairs.sp(g)
```

Arguments

g

graph object with edge weights given

Details

Compute shortest paths between every pair of vertices for a dense graph. It works on both undirected and directed graph. The result is given as a distance matrix. The matrix is symmetric for an undirected graph, and asymmetric (very likely) for a directed graph. For a sparse graph, the johnson.all.pairs.sp functions should be used instead.

See documentation on these algorithms in Boost Graph Library for more details.

Value

A matrix of shortest path lengths between all pairs of nodes in the graph.

Author(s)

Li Long <li.long@isb-sib.ch>

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

The Boost Graph Library: User Guide and Reference Manual; by Jeremy G. Siek, Lie-Quan Lee, and Andrew Lumsdaine; (Addison-Wesley, Pearson Education Inc., 2002), xxiv+321pp. ISBN 0-201-72914-8

See Also

```
johnson.all.pairs.sp
```

```
con <- file(system.file("XML/conn.gxl", package="RBGL"), open="r")
coex <- fromGXL(con)
close(con)
floyd.warshall.all.pairs.sp(coex)</pre>
```

24 gprofile

gprofile

Compute profile for a graph

Description

Compute profile for a graph

Usage

```
gprofile(g)
```

Arguments

g

an instance of the graph class

Details

The profile of a given graph is the sum of bandwidths for all the vertices in the graph.

See documentation on this function in Boost Graph Library for more details.

Value

```
profile the profile of the graph
```

Author(s)

Li Long <li.long@isb-sib.ch>

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

The Boost Graph Library: User Guide and Reference Manual; by Jeremy G. Siek, Lie-Quan Lee, and Andrew Lumsdaine; (Addison-Wesley, Pearson Education Inc., 2002), xxiv+321pp. ISBN 0-201-72914-8

```
con <- file(system.file("XML/dijkex.gxl",package="RBGL"), open="r")
coex <- fromGXL(con)
close(con)
gprofile(coex)</pre>
```

graphGenerator 25

graphGenerator

Generate an undirected graph with adjustable clustering coefficient

Description

Generate an undirected graph with adjustable clustering coefficient

Usage

```
graphGenerator(n, d, o)
```

Arguments

n no. of nodes in the generated graph
 d parameter for preferential attachment
 o parameter for triple generation

Details

The graph generator works according to the prefential attachment rule. It also generates graphs with adjustable clustering coefficient. Parameter d specifies how many preferred edges a new node has. Parameter o limits how many triples to add to a new node.

See reference for details.

Value

Author(s)

Li Long < li.long@isb-sib.ch>

References

Approximating Clustering Coefficient and Transitivity, T. Schank, D. Wagner, Journal of Graph Algorithms and Applications, Vol. 9, No. 2 (2005).

See Also

clusteringCoef, transitivity, clusteringCoefAppr

```
n <- 20
d <- 6
o <- 3
gg <- graphGenerator(n, d, o)</pre>
```

26 highlyConnSG

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Compute highly connected subgraphs for an undirected graph

Description

Compute highly connected subgraphs for an undirected graph

Usage

```
highlyConnSG(q, sat=3, ldv=c(3,2,1))
```

Arguments

g an instance of the graph class with edgemode "undirected"
sat singleton adoption threshold, positive integer
ldv heuristics to remove lower degree vertice, a decreasing sequence of positive

integer

Details

A graph G with n vertices is highly connected if its connectivity k(G) > n/2. The HCS algorithm partitions a given graph into a set of highly connected subgraphs, by using minimum-cut algorithm recursively. To improve performance, the approach is refined by adopting singletons, removing low degree vertices and merging clusters.

On singleton adoption: after each round of partition, some highly connected subgraphs could be singletons (i.e., a subgraph contains only one node). To reduce the number of singletons, therefore reduce number of clusters, we try to get "normal" subgraphs to "adopt" them. If a singleton, s, has n neighbours in a highly connected subgraph c, and n >sat, we add s to c. To adapt to the modified subgraphs, this adoption process is repeated until no further such adoption.

On lower degree vertices: when the graph has low degree vertices, minimum-cut algorithm will just repeatedly separate these vertices from the rest. To reduce such expensive and non-informative computation, we "remove" these low degree vertices first before applying minimum-cut algorithm. Given ldv = (d1, d2, ..., dp), (d[i] > d[i+1] > 0), we repeat the following (i from 1 to p): remove all the highly-connected-subgraph found so far; remove vertices with degrees < di; find highly-connected-subgraphs; perform singleton adoptions.

The Boost implementation does not support self-loops, therefore we signal an error and suggest that users remove self-loops using the function removeSelfLoops function. This change does affect degree, but the original article makes no specific reference to self-loops.

Value

A list of clusters, each is given as vertices in the graph.

Author(s)

Li Long <li.long@isb-sib.ch>

References

A Clustering Algorithm based on Graph Connectivity by E. Hartuv, R. Shamir, 1999.

See Also

```
edgeConnectivity, minCut, removeSelfLoops
```

Examples

```
con <- file(system.file("XML/hcs.gxl",package="RBGL"))
coex <- fromGXL(con)
close(con)
highlyConnSG(coex)</pre>
```

incremental.components

Compute connected components for an undirected graph

Description

Compute connected components for an undirected graph

Usage

```
init.incremental.components(g)
incremental.components(g)
same.component(g, node1, node2)
```

Arguments

g an instance of the graph class node1 one vertex of the given graph node2 another vertex of the given graph

Details

This family of functions work together to calculate the connected components of an undirected graph. The algorithm is based on the disjoint-sets. It works where the graph is growing by adding new edges. Call "init.incremental.components" to initialize the calculation on a new graph. Call "incremental.components" to re-calculate connected components after growing the graph. Call "same.component" to learn if two given vertices are in the same connected components. Currently, the codes can only handle ONE incremental graph at a time. When you start working on another graph by calling "init.incremental.components", the disjoint-sets info on the previous graph is lost. See documentation on Incremental Connected Components in Boost Graph Library for more details.

Value

Output from init.incremental.components is a list of component numbers for each vertex in the graph.

Output from incremental.components is a list of component numbers for each vertex in the graph.

Output from same.component is true if both nodes are in the same connected component, otherwise it's false.

Author(s)

Li Long <li.long@isb-sib.ch>

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

The Boost Graph Library: User Guide and Reference Manual; by Jeremy G. Siek, Lie-Quan Lee, and Andrew Lumsdaine; (Addison-Wesley, Pearson Education Inc., 2002), xxiv+321pp. ISBN 0-201-72914-8

See Also

connComp, connectedComp, strongComp

Examples

```
con <- file(system.file("XML/conn2.gxl",package="RBGL"), open="r")
coex <- fromGXL(con)
close(con)
init.incremental.components(coex)
incremental.components(coex)
v1 <- 1
v2 <- 5
same.component(coex, v1, v2)</pre>
```

isKuratowskiSubgraph

isKuratowskiSubgraph

Description

isKuratowskiSubgraph description

Usage

```
isKuratowskiSubgraph(g)
```

Arguments

g

instance of class graphNEL from Bioconductor graph class

Author(s)

Li Long <li.long@isb-sib.ch>

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

The Boost Graph Library: User Guide and Reference Manual; by Jeremy G. Siek, Lie-Quan Lee, and Andrew Lumsdaine; (Addison-Wesley, Pearson Education Inc., 2002), xxiv+321pp. ISBN 0-201-72914-8

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Examples

```
V <- LETTERS[1:6]
g <- new("graphNEL", nodes=V, edgemode="undirected")</pre>
g \leftarrow addEdge(V[1+0], V[1+1], g)
g \leftarrow addEdge(V[1+0], V[2+1], g)
g \leftarrow addEdge(V[1+0], V[3+1], g)
g \leftarrow addEdge(V[1+0], V[4+1], g)
g \leftarrow addEdge(V[1+0], V[5+1], g)
g \leftarrow addEdge(V[1+1], V[2+1], g)
g \leftarrow addEdge(V[1+1], V[3+1], g)
g \leftarrow addEdge(V[1+1], V[4+1], g)
g \leftarrow addEdge(V[1+1], V[5+1], g)
g \leftarrow addEdge(V[1+2], V[3+1], g)
g \leftarrow addEdge(V[1+2], V[4+1], g)
g \leftarrow addEdge(V[1+2], V[5+1], g)
g \leftarrow addEdge(V[1+3], V[4+1], g)
g \leftarrow addEdge(V[1+3], V[5+1], g)
g \leftarrow addEdge(V[1+4], V[5+1], g)
x4 <- isKuratowskiSubgraph(g)
x4
```

isomorphism

Compute isomorphism from vertices in one graph to those in another graph

Description

Compute isomorphism from vertices in one graph to those in another graph

Usage

```
isomorphism(g1, g2)
```

Arguments

```
g1 one instance of the graph class
g2 one instance of the graph class
```

Details

As stated in documentation on isomorphism in Boost Graph Library: An isomorphism is a 1-to-1 mapping of the vertices in one graph to the vertices of another graph such that adjacency is preserved. Another words, given graphs G1 = (V1,E1) and G2 = (V2,E2) an isomorphism is a function f such that for all pairs of vertices a,b in V1, edge (a,b) is in E1 if and only if edge (f(a),f(b)) is in E2.

Value

Output is true if there exists an isomorphism between g1 and g2, otherwise it's false.

Author(s)

Li Long <li.long@isb-sib.ch>

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

The Boost Graph Library: User Guide and Reference Manual; by Jeremy G. Siek, Lie-Quan Lee, and Andrew Lumsdaine; (Addison-Wesley, Pearson Education Inc., 2002), xxiv+321pp. ISBN 0-201-72914-8

Examples

```
con1 <- file(system.file("XML/dijkex.gxl",package="RBGL"), open="r")
g1 <- fromGXL(con1)
close(con1)

con2 <- file(system.file("XML/conn2.gxl",package="RBGL"), open="r")
g2 <- fromGXL(con2)
close(con2)
isomorphism(g1, g2)</pre>
```

isStraightLineDrawing

isStraightLineDrawing

Description

isStraightLineDrawing description

Usage

```
isStraightLineDrawing(g, drawing)
```

Arguments

g instance of class graphNEL from Bioconductor graph class

drawing coordinates of node positions

Author(s)

Li Long <li.long@isb-sib.ch>

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

The Boost Graph Library: User Guide and Reference Manual; by Jeremy G. Siek, Lie-Quan Lee, and Andrew Lumsdaine; (Addison-Wesley, Pearson Education Inc., 2002), xxiv+321pp. ISBN 0-201-72914-8

is.triangulated 31

Examples

```
V <- LETTERS[1:7]
g <- new("graphNEL", nodes=V, edgemode="undirected")</pre>
g <- addEdge(V[1+0], V[1+1], g)
g \leftarrow addEdge(V[1+1], V[2+1], g)
g \leftarrow addEdge(V[1+2], V[3+1], g)
g \leftarrow addEdge(V[1+3], V[0+1], g)
g \leftarrow addEdge(V[1+3], V[4+1], g)
g \leftarrow addEdge(V[1+4], V[5+1], g)
g \leftarrow addEdge(V[1+5], V[6+1], g)
g \leftarrow addEdge(V[1+6], V[3+1], g)
g \leftarrow addEdge(V[1+0], V[4+1], g)
g \leftarrow addEdge(V[1+1], V[3+1], g)
g \leftarrow addEdge(V[1+3], V[5+1], g)
g \leftarrow addEdge(V[1+2], V[6+1], g)
g \leftarrow addEdge(V[1+1], V[4+1], g)
g \leftarrow addEdge(V[1+1], V[5+1], g)
g \leftarrow addEdge(V[1+1], V[6+1], g)
x3 <- chrobakPayneStraightLineDrawing(g)
x8 <- isStraightLineDrawing(g, x3)
x8
```

is.triangulated

Decide if a graph is triangulated

Description

Decide if a graph is triangulated

Usage

```
is.triangulated(g)
```

Arguments

g

an instance of the graph class

Details

An undirected graph G = (V, E) is triangulated (i.e. chordal) if all cycles [v1, v2, ..., vk] of length 4 or more have a chord, i.e., an edge [vi, vj] with $j! = i + /-1 \pmod{k}$

An equivalent definition of chordal graphs is:

G is chordal iff either G is an empty graph, or there is an v in V such that

- 1. the neighborhood of v (i.e., v and its adjacent nodes) forms a clique, and
- 2. recursively, G-v is chordal

johnson.all.pairs.sp

Value

The return value is TRUE if g is triangulated and FALSE otherwise. An error is thrown if the graph is not undirected; you might use ugraph to compute the underlying graph.

Author(s)

Li Long <li.long@isb-sib.ch>

References

Combinatorial Optimization: algorithms and complexity (p. 403) by C. H. Papadimitriou, K. Steiglitz

Examples

```
con1 <- file(system.file("XML/conn.gxl",package="RBGL"), open="r")
coex <- fromGXL(con1)
close(con1)

is.triangulated(coex)

con2 <- file(system.file("XML/hcs.gxl",package="RBGL"), open="r")
coex <- fromGXL(con2)
close(con2)

is.triangulated(coex)</pre>
```

```
johnson.all.pairs.sp
```

compute shortest path distance matrix for all pairs of nodes

Description

compute shortest path distance matrix for all pairs of nodes

Usage

```
johnson.all.pairs.sp(g)
```

Arguments

g

graph object for which edgeMatrix and edgeWeights are defined

Details

Uses BGL algorithm.

Value

matrix of shortest path lengths, read from row node to col node

kCliques 33

Author(s)

Vince Carey <stvjc@channing.harvard.edu>

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

The Boost Graph Library: User Guide and Reference Manual; by Jeremy G. Siek, Lie-Quan Lee, and Andrew Lumsdaine; (Addison-Wesley, Pearson Education Inc., 2002), xxiv+321pp. ISBN 0-201-72914-8

See Also

```
bellman.ford.sp, dag.sp, dijkstra.sp, sp.between
```

Examples

```
con <- file(system.file("dot/joh.gxl", package="RBGL"), open="r")
z <- fromGXL(con)
close(con)
johnson.all.pairs.sp(z)</pre>
```

kCliques

Find all the k-cliques in an undirected graph

Description

Find all the k-cliques in an undirected graph

Usage

```
kCliques(g)
```

Arguments

g

an instance of the graph class

Details

Notice that there are different definitions of k-clique in different context.

In computer science, a k-clique of a graph is a clique, i.e., a complete subgraph, of k nodes.

In Social Network Analysis, a k-clique in a graph is a subgraph where the distance between any two nodes is no greater than k.

Here we take the definition in Social Network Analysis.

Let D be a matrix, D[i][j] is the shortest path from node i to node j. Algorithm is outlined as following: (1) use Johnson's algorithm to fill D; let $N = \max(D[i][j])$ for all i, j; (2) each edge is a 1-clique by itself; (3) for k = 2, ..., N, try to expand each (k-1)-clique to k-clique: (3.1) consider a (k-1)-clique the current k-clique KC; (3.2) repeat the following: if for all nodes j in KC, $D[v][j] \le k$, add node v to KC; (3.3) eliminate duplicates; (4) the whole graph is N-clique.

kCores

Value

A list of length N; k-th entry (k = 1, ..., N) is a list of all the k-cliques in graph g.

Author(s)

Li Long <li.long@isb-sib.ch>

References

Social Network Analysis: Methods and Applications. By S. Wasserman and K. Faust, pp. 258.

Examples

```
con <- file(system.file("XML/snacliqueex.gxl",package="RBGL"))
coex <- fromGXL(con)
close(con)
kCliques(coex)</pre>
```

kCores

Find all the k-cores in a graph

Description

Find all the k-cores in a graph

Usage

```
kCores(g, EdgeType=c("in", "out"))
```

Arguments

g an instance of the graph class

EdgeType what types of edges to be considered when g is directed

Details

A k-core in a graph is a subgraph where each node is adjacent to at least a minimum number, k, of the other nodes in the subgraph.

A k-core in a graph may not be connected.

The core number for each node is the highest k-core this node is in. A node in a k-core will be, by definition, in a (k-1)-core.

The implementation is based on the algorithm by V. Batagelj and M. Zaversnik, 2002.

The example snacoreex.gxl is in the paper by V. Batagelj and M. Zaversnik, 2002.

Value

A vector of the core numbers for all the nodes in g.

lambdaSets 35

Author(s)

Li Long <li.long@isb-sib.ch>

References

Social Network Analysis: Methods and Applications. By S. Wasserman and K. Faust, pp. 266. An O(m) Algorithm for Cores decomposition of networks, by V. Batagelj and M. Zaversnik, 2002.

Examples

```
con1 <- file(system.file("XML/snacoreex.gxl",package="RBGL"))
kcoex <- fromGXL(con1)
close(con1)

kCores(kcoex)

con2 <- file(system.file("XML/conn2.gxl",package="RBGL"))
kcoex2 <- fromGXL(con2)
close(con2)

kCores(kcoex2)
kCores(kcoex2, "in")
kCores(kcoex2, "out")</pre>
```

lambdaSets

Find all the lambda-sets in an undirected graph

Description

Find all the lambda-sets in an undirected graph

Usage

```
lambdaSets(q)
```

Arguments

g

an instance of the graph class

Details

From reference (1), p. 270: A set of nodes is a lambda-set if any pair of nodes in the lambda set has larger edge connectivity than any pair of nodes consisting of one node from within the lambda set and a second node from outside the lambda set.

As stated in reference (2), a lambda set is a maximal subset of nodes who have more edge-independent paths connecting them to each other than to outsiders.

A lambda set could be characterized by the minimum edge connectivity k among its members, and could be called lambda-k sets.

Let N be maximum edge connectivity of graph g, we output all the lambda-k set for all k = 1, ..., N.

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Value

Maximum edge connectivity, N, of the graph g, and A list of length N; k-th entry (k = 1, ..., N) is a list of all the lambda-k sets in graph g.

Author(s)

Li Long <li.long@isb-sib.ch>

References

- (1) Social Network Analysis: Methods and Applications. By S. Wasserman and K. Faust, pp. 269.
- (2) LS sets, lambda sets and other cohesive subsets. By S. P. Borgatti, M. G. Everett, P. R. Shirey, Social Networks 12 (1990) p. 337-357

Examples

```
con <- file(system.file("XML/snalambdaex.gxl",package="RBGL"))
coex <- fromGXL(con)
close(con)
lambdaSets(coex)</pre>
```

layout

Layout an undirected graph in 2D

Description

Layout an undirected graph in 2D

Usage

```
circleLayout(g, radius=1)
kamadaKawaiSpringLayout( g, edge_or_side=1, es_length=1 )
fruchtermanReingoldForceDirectedLayout(g, width=1, height=1)
randomGraphLayout(g, minX=0, maxX=1, minY=0, maxY=1)
```

Arguments

g	an instance of the graph class with edgemode "undirected"
radius	radius of a regular n-polygon
edge_or_side	boolean indicating the length is for an edge or for a side, default is for an edge
es_length	the length of an edge or a side for layout
width	the width of the dislay area, all x coordinates fall in [-width/2, width/2]
height	the height of the display area, all y coordinates fall in [-height/2, height/2]
minX	minimum x coordinate
maxX	maximum x coordinate
minY	minimum y coordinate
maxY	maximum y coordinate

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Details

If you want to simply draw a graph, you should consider using package Rgraphviz. The layout options in package Rgraphviz: neato, circo and fdp, correspond to kamadaKawaiSpringLayout, circleLayout and fruchtermanReingoldForceDirectedLayout, respectively.

Function circleLayout layouts the graph with the vertices at the points of a regular n-polygon. The distance from the center of the polygon to each point is determined by the radius parameter.

Function kamadaKawaiSpringLayout provides Kamada-Kawai spring layout for connected, undirected graphs. User provides either the unit length e of an edge in the layout or the length of a side s of the display area.

Function randomGraphLayout places the points of the graph at random locations.

Function fruchtermanReingoldForceDirectedLayout performs layout of unweighted, undirected graphs. It's a force-directed algorithm. The BGL implementation doesn't handle disconnected graphs very well, since it doesn't explicitly give each connected component a region proportional to its size.

See documentation on this function in Boost Graph Library for more details.

Value

A $(2 \times n)$ matrix, where n is the number of nodes in the graph, each column gives the (x, y)-coordinates for the corresponding node.

Author(s)

Li Long <li.long@isb-sib.ch>

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

The Boost Graph Library: User Guide and Reference Manual; by Jeremy G. Siek, Lie-Quan Lee, and Andrew Lumsdaine; (Addison-Wesley, Pearson Education Inc., 2002), xxiv+321pp. ISBN 0-201-72914-8

See Also

```
layoutGraph
```

```
con <- file(system.file("XML/conn.gxl",package="RBGL"), open="r")
coex <- fromGXL(con)
close(con)

coex <- ugraph(coex)

circleLayout(coex)

kamadaKawaiSpringLayout(coex)

randomGraphLayout(coex)

fruchtermanReingoldForceDirectedLayout(coex, 10, 10)</pre>
```

38 makeBiconnectedPlanar

```
makeBiconnectedPlanar
```

makeBiconnectedPlanar

Description

makeBiconnectedPlanar description

Usage

```
makeBiconnectedPlanar(g)
```

Arguments

g

instance of class graphNEL from Bioconductor graph class

Author(s)

Li Long <li.long@isb-sib.ch>

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

The Boost Graph Library: User Guide and Reference Manual; by Jeremy G. Siek, Lie-Quan Lee, and Andrew Lumsdaine; (Addison-Wesley, Pearson Education Inc., 2002), xxiv+321pp. ISBN 0-201-72914-8

```
V <- LETTERS[1:11]</pre>
g <- new("graphNEL", nodes=V, edgemode="undirected")</pre>
g \leftarrow addEdge(V[1+0], V[1+1], g)
g \leftarrow addEdge(V[1+2], V[3+1], g)
g \leftarrow addEdge(V[1+3], V[0+1], g)
g \leftarrow addEdge(V[1+3], V[4+1], g)
g \leftarrow addEdge(V[1+4], V[5+1], g)
g \leftarrow addEdge(V[1+5], V[3+1], g)
g \leftarrow addEdge(V[1+5], V[6+1], g)
g \leftarrow addEdge(V[1+6], V[7+1], g)
g \leftarrow addEdge(V[1+7], V[8+1], g)
g \leftarrow addEdge(V[1+8], V[5+1], g)
g \leftarrow addEdge(V[1+8], V[9+1], g)
g \leftarrow addEdge(V[1+0], V[10+1], g)
x6 <- makeBiconnectedPlanar(g)
х6
```

makeConnected 39

makeConnected

make Connected

Description

makeConnected description

Usage

```
makeConnected(g)
```

Arguments

g

instance of class graphNEL from Bioconductor graph class

Author(s)

Li Long <li.long@isb-sib.ch>

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

The Boost Graph Library: User Guide and Reference Manual; by Jeremy G. Siek, Lie-Quan Lee, and Andrew Lumsdaine; (Addison-Wesley, Pearson Education Inc., 2002), xxiv+321pp. ISBN 0-201-72914-8

```
V <- LETTERS[1:11]
g <- new("graphNEL", nodes=V, edgemode="undirected")
g <- addEdge(V[1+0], V[1+1], g)
g <- addEdge(V[1+2], V[3+1], g)
g <- addEdge(V[1+3], V[4+1], g)
g <- addEdge(V[1+5], V[6+1], g)
g <- addEdge(V[1+6], V[7+1], g)
g <- addEdge(V[1+8], V[9+1], g)
g <- addEdge(V[1+9], V[10+1], g)
g <- addEdge(V[1+9], V[8+1], g)</pre>
x5 <- makeConnected(g)
x5</pre>
```

40 makeMaximalPlanar

```
make Maximal Planar make Maximal Planar
```

Description

makeMaximalPlanar description

Usage

```
makeMaximalPlanar(g)
```

Arguments

g

instance of class graphNEL from Bioconductor graph class

Author(s)

Li Long <li.long@isb-sib.ch>

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

The Boost Graph Library: User Guide and Reference Manual; by Jeremy G. Siek, Lie-Quan Lee, and Andrew Lumsdaine; (Addison-Wesley, Pearson Education Inc., 2002), xxiv+321pp. ISBN 0-201-72914-8

```
V <- LETTERS[1:10]
g <- new("graphNEL", nodes=V, edgemode="undirected")
g <- addEdge(V[1+0], V[1+1], g)
g <- addEdge(V[1+1], V[2+1], g)
g <- addEdge(V[1+2], V[3+1], g)
g <- addEdge(V[1+3], V[4+1], g)
g <- addEdge(V[1+4], V[5+1], g)
g <- addEdge(V[1+5], V[6+1], g)
g <- addEdge(V[1+6], V[7+1], g)
g <- addEdge(V[1+7], V[8+1], g)
g <- addEdge(V[1+7], V[8+1], g)</pre>
x7 <- makeMaximalPlanar(g)
x7</pre>
```

maxClique 41

maxClique

Find all the cliques in a graph

Description

Find all the cliques in a graph

Usage

```
maxClique(g)
```

Arguments

g

an instance of the graph class

Details

Notice the maximum clique problem is NP-complete, which means it cannot be solved by any known polynomial algorithm.

We implemented the algorithm by C. Bron and J. Kerbosch,

Value

```
maxClique list of all cliques in g
```

Author(s)

Li Long < li.long@isb-sib.ch>

References

Finding all cliques of an undirected graph, by C. Bron and J. Kerbosch, Communication of ACM, Sept 1973, Vol 16, No. 9.

```
con1 <- file(system.file("XML/conn.gxl",package="RBGL"), open="r")
coex <- fromGXL(con1)
close(con1)

maxClique(coex)

con2 <- file(system.file("XML/hcs.gxl",package="RBGL"), open="r")
coex <- fromGXL(con2)
close(con2)

maxClique(coex)</pre>
```

42 max.flow

max.flow

Compute max flow for a directed graph

Description

Compute max flow for a directed graph

Usage

```
edmonds.karp.max.flow(g, source, sink)
push.relabel.max.flow(g, source, sink)
kolmogorov.max.flow(g, source, sink)
```

Arguments

g an instance of the graph class with edgemode "directed"

source node name (character) or node number (int) for the source of the flow sink node name (character) or node number (int) for the sink of the flow

Details

Given a directed graph G=(V, E) of a single connected component with a vertex source and a vertex sink. Each arc has a positive real valued capacity, currently it's equivalent to the weight of the arc. The flow of the network is the net flow entering the vertex sink. The maximum flow problem is to determine the maximum possible value for the flow to the sink and the corresponding flow values for each arc.

See documentation on these algorithms in Boost Graph Library for more details.

Value

A list of

maxflow the max flow from source to sink

edges the nodes of the arcs with non-zero capacities

flows the flow values of the arcs with non-zero capacities

Author(s)

Li Long <li.long@isb-sib.ch>

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

The Boost Graph Library: User Guide and Reference Manual; by Jeremy G. Siek, Lie-Quan Lee, and Andrew Lumsdaine; (Addison-Wesley, Pearson Education Inc., 2002), xxiv+321pp. ISBN 0-201-72914-8

See Also

```
minCut, edgeConnectivity
```

maximumCycleRatio 43

Examples

```
con <- file(system.file("XML/dijkex.gxl",package="RBGL"), open="r")
g <- fromGXL(con)
close(con)

ans1 <- edmonds.karp.max.flow(g, "B", "D")
ans2 <- edmonds.karp.max.flow(g, 3, 2)  # 3 and 2 equivalent to "C" and "B"

ans3 <- push.relabel.max.flow(g, 2, 4)  # 2 and 4 equivalent to "B" and "D"
ans4 <- push.relabel.max.flow(g, "C", "B")

# error in the following now
#ans5 <- kolmogorov.max.flow(g, "B", "D")
#ans6 <- kolmogorov.max.flow(g, 3, 2)</pre>
```

maximumCycleRatio maximumCycleRatio

Description

maximumCycleRatio description

Usage

```
maximumCycleRatio(g)
```

Arguments

g

instance of class graphNEL from Bioconductor graph class

Author(s)

Li Long < li.long@isb-sib.ch>

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

44 minCut

minCut

Compute min-cut for an undirected graph

Description

Compute min-cut for an undirected graph

Usage

```
minCut(g)
```

Arguments

a

an instance of the graph class with edgemode "undirected"

Details

Given an undirected graph G=(V, E) of a single connected component, a cut is a partition of the set of vertices into two non-empty subsets S and V-S, a cost is the number of edges that are incident on one vertex in S and one vertex in V-S. The min-cut problem is to find a cut (S, V-S) of minimum cost.

For simplicity, the returned subset S is the smaller of the two subsets.

Value

A list of

mincut the number of edges to be severed to obtain the minimum cut

S the smaller subset of vertices in the minimum cut V-S the other subset of vertices in the minimum cut

Author(s)

Li Long <li.long@isb-sib.ch>

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

The Boost Graph Library: User Guide and Reference Manual; by Jeremy G. Siek, Lie-Quan Lee, and Andrew Lumsdaine; (Addison-Wesley, Pearson Education Inc., 2002), xxiv+321pp. ISBN 0-201-72914-8

See Also

```
edgeConnectivity
```

```
con <- file(system.file("XML/conn.gxl",package="RBGL"), open="r")
coex <- fromGXL(con)
close(con)
minCut(coex)</pre>
```

minimumCycleRatio 45

minimumCycleRatio minimumCycleRatio

Description

minimumCycleRatio description

Usage

```
minimumCycleRatio(g)
```

Arguments

g instance of class graphNEL from Bioconductor graph class

Author(s)

Li Long li.long@isb-sib.ch

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

The Boost Graph Library: User Guide and Reference Manual; by Jeremy G. Siek, Lie-Quan Lee, and Andrew Lumsdaine; (Addison-Wesley, Pearson Education Inc., 2002), xxiv+321pp. ISBN 0-201-72914-8

mstree.kruskal

Kruskal's minimum spanning tree in boost

Description

compute the minimum spanning tree (MST) for a graph and return a representation in matrices

Usage

```
mstree.kruskal(x)
```

Arguments

x instance of class graph

Details

calls to kruskal minimum spanning tree algorithm of Boost graph library

46 mstree.prim

Value

a list

edgeList a matrix m of dimension 2 by number of edges in the MST, with m[i,j] the jth

node in edge i

weights a vector of edge weights corresponding to the columns of edgeList

nodes the vector of nodes of the input graph x

Author(s)

VJ Carey <stvjc@channing.harvard.edu>

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

The Boost Graph Library: User Guide and Reference Manual; by Jeremy G. Siek, Lie-Quan Lee, and Andrew Lumsdaine; (Addison-Wesley, Pearson Education Inc., 2002), xxiv+321pp. ISBN 0-201-72914-8

See Also

```
prim.minST
```

Examples

```
con1 <- file(system.file("XML/kmstEx.gxl",package="RBGL"), open="r")
km <- fromGXL(con1)
close(con1)

mstree.kruskal(km)
edgeData(km, "B", "D", "weight") <- 1.1
edgeData(km, "B", "E", "weight") <- .95
mstree.kruskal(km)

con2 <- file(system.file("XML/telenet.gxl",package="RBGL"), open="r")
km2 <- fromGXL(con2)
close(con2)

m <- mstree.kruskal(km2)
print(sum(m[[2]]))</pre>
```

mstree.prim

Compute minimum spanning tree for an undirected graph

Description

Compute minimum spanning tree for an undirected graph

Usage

```
mstree.prim(g)
prim.minST(g)
```

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Arguments

g

an instance of the graph class with edgemode "undirected"

Details

This is Prim's algorithm for solving the minimum spanning tree problem for an undirected graph with weighted edges.

The function prim.minST is depricated and renamed, use mstree.prim instead.

See documentations on this function in Boost Graph Library for more details.

Value

A list of

edges the edges that form the minimum spanning tree
weights the total weight of the minimum spanning tree

Author(s)

Li Long <li.long@isb-sib.ch>

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

The Boost Graph Library: User Guide and Reference Manual; by Jeremy G. Siek, Lie-Quan Lee, and Andrew Lumsdaine; (Addison-Wesley, Pearson Education Inc., 2002), xxiv+321pp. ISBN 0-201-72914-8

See Also

```
mstree.kruskal
```

Examples

```
con <- file(system.file("XML/conn2.gxl",package="RBGL"))
coex <- fromGXL(con)
close(con)

mstree.prim(coex)</pre>
```

Ordering

Compute vertex ordering for an undirected graph

Description

Compute vertex ordering for an undirected graph

Usage

```
cuthill.mckee.ordering(g)
minDegreeOrdering(g, delta=0)
sloan.ordering(g, w1=1, w2=2)
```

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Arguments

g	an instance of the graph class with edgemode "undirected"
delta	Multiple elimination control variable. If it is larger than or equal to zero then multiple elimination is enabled. The value of delta specifies the difference between the minimum degree and the degree of vertices that are to be eliminated.
w1	First heuristic weight for the Sloan algorithm.
w2	Second heuristic weight for the Sloan algorithm.

Details

The following details were obtained from the documentation of these algorithms in Boost Graph Library and readers are referred their for even more detail. The goal of the Cuthill-Mckee (and reverse Cuthill-Mckee) ordering algorithm is to reduce the bandwidth of a graph by reordering the indices assigned to each vertex.

The minimum degree ordering algorithm is a fill-in reduction matrix reordering algorithm.

The goal of the Sloan ordering algorithm is to reduce the profile and the wavefront of a graph by reordering the indices assigned to each vertex.

The goal of the King ordering algorithm is to reduce the bandwidth of a graph by reordering the indices assigned to each vertex.

Value

rms.wavefront

```
cuthill.mckee.ordering
                 returns a list with elements:
reverse cuthill.mckee.ordering
                 the vertices in the new ordering
original bandwidth
                 bandwidth before reordering vertices
new bandwidth
                 bandwidth after reordering of vertices
minDegreeOrdering
                 return a list with elements:
inverse_permutation
                 the new vertex ordering, given as the mapping from the new indices to the old
                 indices
permutation
                 the new vertex ordering, given as the mapping from the old indices to the new
                 indices
sloan.ordering
                 returns a list with elements:
sloan.ordering
                 the vertices in the new ordering
                 bandwidth of the graph after reordering
bandwidth
                 profile of the graph after reordering
profile
maxWavefront maxWavefront of the graph after reordering
aver.wavefront
```

aver.wavefront of the graph after reordering

rms.wavefront of the graph after reordering

Author(s)

Li Long li.long@isb-sib.ch

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

The Boost Graph Library: User Guide and Reference Manual; by Jeremy G. Siek, Lie-Quan Lee, and Andrew Lumsdaine; (Addison-Wesley, Pearson Education Inc., 2002), xxiv+321pp. ISBN 0-201-72914-8

Examples

```
con <- file(system.file("XML/dijkex.gxl",package="RBGL"), open="r")
coex <- fromGXL(con)
close(con)

coex <- ugraph(coex)
cuthill.mckee.ordering(coex)
minDegreeOrdering(coex)
sloan.ordering(coex)</pre>
```

planarCanonicalOrdering

planarCanonicalOrdering

Description

planarCanonicalOrdering description

Usage

```
planarCanonicalOrdering(g)
```

Arguments

g

instance of class graphNEL from Bioconductor graph class

Author(s)

Li Long li.long@isb-sib.ch>

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

50 planarFaceTraversal

Examples

```
V <- LETTERS[1:6]
g <- new("graphNEL", nodes=V, edgemode="undirected")</pre>
g \leftarrow addEdge(V[1+0], V[1+1], g)
g \leftarrow addEdge(V[1+1], V[2+1], g)
g \leftarrow addEdge(V[1+2], V[3+1], g)
g \leftarrow addEdge(V[1+3], V[4+1], g)
g \leftarrow addEdge(V[1+4], V[5+1], g)
g \leftarrow addEdge(V[1+5], V[0+1], g)
g \leftarrow addEdge(V[1+0], V[2+1], g)
g \leftarrow addEdge(V[1+0], V[3+1], g)
g \leftarrow addEdge(V[1+0], V[4+1], g)
g \leftarrow addEdge(V[1+1], V[3+1], g)
g <- addEdge(V[1+1], V[4+1], g)
g \leftarrow addEdge(V[1+1], V[5+1], g)
x2 <- planarCanonicalOrdering(g)
x2
```

```
planarFaceTraversal
```

planarFaceTraversal

Description

planarFaceTraversal description

Usage

```
planarFaceTraversal(g)
```

Arguments

g instance of class graphNEL from Bioconductor graph class

Author(s)

Li Long <li.long@isb-sib.ch>

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

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Examples

```
V <- LETTERS[1:9]
g <- new("graphNEL", nodes=V, edgemode="undirected")</pre>
g \leftarrow addEdge(V[1+0],V[1+1],g)
g <- addEdge(V[1+1], V[1+2], g)</pre>
g \leftarrow addEdge(V[1+3],V[1+4],g)
g <- addEdge(V[1+4],V[1+5],g)</pre>
g \leftarrow addEdge(V[1+6],V[1+7],g)
g \leftarrow addEdge(V[1+7],V[1+8],g)
g \leftarrow addEdge(V[1+0],V[1+3],g)
g \leftarrow addEdge(V[1+3],V[1+6],g)
g <- addEdge(V[1+1],V[1+4],g)</pre>
g <- addEdge(V[1+4],V[1+7],g)</pre>
g \leftarrow addEdge(V[1+2],V[1+5],g)
g <- addEdge(V[1+5],V[1+8],g)</pre>
x1 <- planarFaceTraversal(g)</pre>
x1
```

RBGL.overview

RBGL.overview

Description

The RBGL package consists of a number of interfaces to the Boost C++ library for graph algorithms. This page follows, approximately, the chapter structure of the monograph on the Boost Graph Library by Siek et al., and gives hyperlinks to documentation on R functions currently available, along with the names of formal parameters to these functions.

basicAlgs

Functions	parameters
bandwidth	g
bfs	object, node, checkConn
dfs	object, node, checkConn
edgeConnectivity	g
gprofile	g
isomorphism	g1,g2
minCut	g
transitive.closure	g
tsort	X

ShortestPaths

Functions

parameters

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```
bellman.ford.sp g,start
dag.sp g,start
dijkstra.sp g,start
extractPath s,f,pens
johnson.all.pairs.sp g
sp.between g,start,finish
sp.between.old g,start,finish
sp.between.scalar g,start,finish
```

MinimumSpanningTree

Functions	parameters
mstree.kruskal	X
prim.minST	g

Connected Components

Functions	parameters
connectedComp	g
highlyConnSG	g,sat,ldv
incremental.components	g
<pre>init.incremental.components</pre>	g
same.component	g,node1,node2
strongComp	g

MaximumFlow

Functions	parameters
edmonds.karp.max.flow	g,source,sink
<pre>push.relabel.max.flow</pre>	g,source,sink

SparseMatrixOrdering

Functions	parameters
cuthill.mckee.ordering	g
minDegreeOrdering	g,delta
sloan.ordering	g,w1,w2

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LayoutAlgorithms

```
Functions parameters
circle.layout g,radius
kamada.kawai.spring.layout g,edge_or_side,es_length
```

GraphClustering

```
Functions parameters
betweenness.centrality.clustering g,threshold,normalize
```

Betweenness

```
Functions parameters brandes.betweenness.centrality g
```

Wavefront

Functions	parameters
aver.wavefront	g
ith.wavefront	g,start
maxWavefront	g
rms.wavefront	g

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

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 ${\tt removeSelfLoops}$

remove self loops in a graph

Description

remove self loops in a graph

Usage

```
removeSelfLoops(g)
```

Arguments

g

one instance of the graph class

Details

If a given graph contains self-loop(s), removeSelfLoops removes them. This is for those functions that cannot handle graphs with self-loops.

Value

A new graph without self loops.

Author(s)

Li Long < li.long@isb-sib.ch>

Examples

```
con <- file(system.file("XML/dijkex.gxl",package="RBGL"))
g1 <- fromGXL(con)
close(con)

g2 <- ugraph(g1)
removeSelfLoops(g2)</pre>
```

separates

A function to test whether a subset of nodes separates two other subsets of nodes.

Description

The function tests to see whether a set of nodes, S1, separates all nodes in a from all nodes in b.

Usage

```
separates(a, b, S1, g)
```

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Arguments

a	The names of the nodes in the from set.
b	The names of the nodes in the to set.
S1	The names of the nodes in the separation set.
g	An instance of the graph class. All nodes named in the other arguments must be nodes of this graph.

Details

The algorithm is quite simple. A subgraph is created by removing the nodes named in S1 from g. Then all paths between elements of a to elements of b are tested for. If any path exists the function returns FALSE, otherwise it returns TRUE.

Value

Either TRUE or FALSE depending on whether S1 separates a from b in g1.

Author(s)

R. Gentleman

References

S. Lauritzen, Graphical Models, OUP.

See Also

```
johnson.all.pairs.sp
```

Examples

```
con <- file(system.file("XML/kmstEx.gxl",package="RBGL"))
km <- fromGXL(con)
close(con)
separates("B", "A", "E", km)
separates("B", "A", "C", km)</pre>
```

```
sloanStartEndVertices
```

sloanStartEndVertices

Description

sloanStartEndVertices description

Usage

```
sloanStartEndVertices(g)
```

56 sp.between

Arguments

instance of class graphNEL from Bioconductor graph class

Author(s)

Li Long <li.long@isb-sib.ch>

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

The Boost Graph Library: User Guide and Reference Manual; by Jeremy G. Siek, Lie-Quan Lee, and Andrew Lumsdaine; (Addison-Wesley, Pearson Education Inc., 2002), xxiv+321pp. ISBN 0-201-72914-8

sp.between

Dijkstra's shortest paths using boost C++

Description

dijkstra's shortest paths

Usage

```
sp.between(g, start, finish, detail=TRUE)
```

Arguments

g instance of class graph

start node name(s) for start of path(s)
finish node name(s) for end of path(s)

detail if TRUE, output additional info on the shortest path

Details

These functions are interfaces to the Boost graph library C++ routines for Dijkstra's shortest paths. Function sp.between.scalar is obsolete.

Value

When start and/or finish are vectors, we use the normal cycling rule in R to match both vectors and try to find the shortest path for each pair.

Function sp.between returns a list of info on the shortest paths. Each such shortest path is designated by its starting node and its ending node. Each element in the returned list contains:

length total length (using edge weights) of this shortest path

,

path_detail if requested, a vector of names of the nodes on the shortest path

,

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```
length_detail
```

if requested, a list of edge weights of this shortest path

.

See pathWeights for caveats about undirected graph representation.

Author(s)

VJ Carey <stvjc@channing.harvard.edu>, Li Long <li.long@isb-sib.ch>

See Also

```
bellman.ford.sp, dag.sp, dijkstra.sp, johnson.all.pairs.sp
```

Examples

```
con <- file(system.file("XML/ospf.gxl",package="RBGL"), open="r")
ospf <- fromGXL(con)
close(con)

dijkstra.sp(ospf,nodes(ospf)[6])

sp.between(ospf, "RT6", "RT1")

sp.between(ospf, c("RT6", "RT2"), "RT1", detail=FALSE)

sp.between(ospf, c("RT6", "RT2"), c("RT1","RT5"))

# see NAs for query on nonexistent path
sp.between(ospf,"N10", "N13")</pre>
```

strongComp

Identify Strongly Connected Components

Description

The strongly connected components in a directed graph are identified and returned as a list.

Usage

```
strongComp(g)
```

Arguments

a

graph with edgemode "directed".

Details

Tarjan's algorithm is used to determine all strongly connected components of a directed graph.

58 transitive.closure

Value

A list whose length is the number of strongly connected components in g. Each element of the list is a vector of the node labels for the nodes in that component.

Author(s)

Vince Carey <stvjc@channing.harvard.edu>

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

The Boost Graph Library: User Guide and Reference Manual; by Jeremy G. Siek, Lie-Quan Lee, and Andrew Lumsdaine; (Addison-Wesley, Pearson Education Inc., 2002), xxiv+321pp. ISBN 0-201-72914-8

See Also

connComp,connectedComp, same.component

Examples

```
con <- file(system.file("XML/kmstEx.gxl",package="RBGL"), open="r")
km <- fromGXL(con)
close(con)

km<- graph::addNode(c("F","G","H"), km)
km<- addEdge("G", "H", km, 1)
km<- addEdge("H", "G", km, 1)
strongComp(km)
connectedComp(ugraph(km))</pre>
```

transitive.closure Compute transitive closure of a directed graph

Description

Compute transitive closure of a directed graph

Usage

```
transitive.closure(g)
```

Arguments

g an instance of the graph class

Details

This function calculates the transitive closure of a directed graph. See documentation on this function in Boost Graph Library for more details.

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Value

An object of class graphNEL.

Author(s)

Li Long < li.long@isb-sib.ch>

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

The Boost Graph Library: User Guide and Reference Manual; by Jeremy G. Siek, Lie-Quan Lee, and Andrew Lumsdaine; (Addison-Wesley, Pearson Education Inc., 2002), xxiv+321pp. ISBN 0-201-72914-8

Examples

```
con <- file(system.file("XML/dijkex.gxl",package="RBGL"))
coex <- fromGXL(con)
close(con)
transitive.closure(coex)</pre>
```

transitivity

Calculate transitivity for an undirected graph

Description

Calculate transitivity for an undirected graph

Usage

```
transitivity(g)
```

Arguments

g

an instance of the graph class

Details

For an undirected graph G, let delta(v) be the number of triangles with v as a node, let tau(v) be the number of triples, i.e., paths of length 2 with v as the center node.

Define transitivity T(G) = sum(delta(v)) / sum(tau(v)), for all v in V.

Value

Transitivity for graph g.

Author(s)

```
LiLong <li.long@isb-sib.ch>
```

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References

Approximating Clustering Coefficient and Transitivity, T. Schank, D. Wagner, Journal of Graph Algorithms and Applications, Vol. 9, No. 2 (2005).

See Also

clusteringCoef, clusteringCoefAppr, graphGenerator

Examples

```
con <- file(system.file("XML/conn.gxl",package="RBGL"))
g <- fromGXL(con)
close(con)

tc <- transitivity(g)</pre>
```

tsort

topological sort of vertices of a digraph

Description

returns vector of zero-based indices of vertices of a DAG in topological sort order

Usage

```
tsort(x) # now x assumed to be Bioconductor graph graphNEL
```

Arguments

Х

instance of class graphNEL from Bioconductor graph class

Details

calls to the topological\sort algorithm of BGL. will check in BGL whether the input is a DAG and return a vector of zeroes (of length length(nodes(x))) if it is not. Thus this function can be used to check for cycles in a digraph.

Value

a numerical vector enumerating vertices in the topological sort sequence, 0-based

Author(s)

VJ Carey <stvjc@channing.harvard.edu>

References

Boost Graph Library (www.boost.org/libs/graph/doc/index.html)

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Examples

```
data(FileDep)
tsind <- tsort(FileDep)
tsind
FD2 <- FileDep
# now introduce a cycle
FD2 <- addEdge("bar_o", "dax_h", FD2, 1)
tsort(FD2)</pre>
```

wavefront

Compute the i-th/max/average/rms wavefront for a graph

Description

Compute the i-th/max/average/rms wavefront for a graph

Usage

```
ith.wavefront(g, start)
maxWavefront(g)
aver.wavefront(g)
rms.wavefront(g)
```

Arguments

```
start a vertex of the graph class
g an instance of the graph class
```

Details

Assorted functions on wavefront of a graph.

Value

Author(s)

Li Long <li.long@isb-sib.ch>

References

```
Boost Graph Library ( www.boost.org/libs/graph/doc/index.html )
```

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See Also

```
edgeConnectivity
```

```
con <- file(system.file("XML/dijkex.gxl",package="RBGL"), open="r")
coex <- fromGXL(con)
close(con)

ss <- 1
ith.wavefront(coex, ss)
maxWavefront(coex)
aver.wavefront(coex)
rms.wavefront(coex)</pre>
```

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